

Representation theory of the monoid of all partial functions on a set and other Ehresmann semigroups

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Shamoon College of Engineering

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Joint work with Stuart Margolis
Bar-Ilan University

Semigroup algebras

- S - finite semigroup.
- \mathbb{k} - field.
- $\mathbb{k}S$ - semigroup algebra.

$$\mathbb{k}S = \left\{ \sum k_i s_i \mid k_i \in \mathbb{k} \quad s_i \in S \right\}$$

$\mathbb{k}S$ is an associative, and finite dimensional \mathbb{k} -algebra.
In this talk it will be also unital.

- For this talk $\mathbb{k} = \mathbb{C}$.
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Simple & Projective modules

- A - associative, unital, finite dimensional \mathbb{k} -algebra.
- **Simple module** - a module with no proper submodules.
 - » Finite up to isomorphism.
- **Ind. module** - cannot be written as direct sum of non-zero modules ($P = P_1 \oplus P_2$)
- **Projective module** - P is projective if $\text{Hom}(P, -)$ is an exact functor.
- **Equivalently:** P is a direct summand of A^k .
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 - ① $S = \mathcal{PT}_n$. The monoid of partial functions on $\{1, \dots, n\}$.
 - ② Generalization: S is a certain kind of Ehresmann semigroup.

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Partial functions

Example

$f \in \mathcal{PT}_4$ given by

$$f = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & \emptyset & 2 & 3 \end{pmatrix}$$

$(f(1) = f(3) = 2, \quad f(4) = 3, \quad f(2) \text{ undefined})$

$$\text{dom}(f) = \{1, 3, 4\}, \quad \text{im}(f) = \{2, 3\}, \quad \text{rank}(f) = |\text{im}(f)| = 2$$

Definition

The kernel of f is an equivalence relation on $\text{dom}(f)$ defined by
 $a \sim b \iff f(a) = f(b)$.

In the above example the kernel classes are $\{1, 3\}, \{4\}$

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Construction of simple modules of \mathcal{PT}_n

- Partial functions $f, g \in \mathcal{PT}_n$ are \mathcal{L} -related $\iff \text{dom}(f) = \text{dom}(g)$ and $\text{ker}(f) = \text{ker}(g)$.
- $f \mathcal{L} g \implies \text{rank}(f) = \text{rank}(g)$. So for \mathcal{L} -class L we can define $\text{rank}(L)$.
- Every \mathcal{L} class contains an idempotent $e \in L$.
- Let L be an \mathcal{L} -class of rank k and let $e \in L$. G_e - The maximal subgroup at e (with unit element e) is isomorphic to S_k (the symmetric group of degree k).

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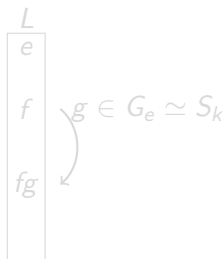
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Natural actions on an \mathcal{L} -class

Let L be an \mathcal{L} -class of rank k and let $e \in L$ be an idempotent. Then $G_e \simeq S_k$ acts on the right of L by right multiplication

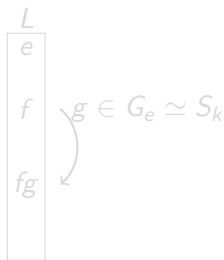
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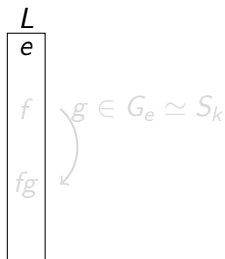
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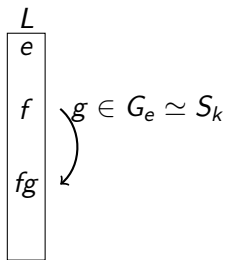
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Let L be an \mathcal{L} -class of rank k and let $e \in L$ be an idempotent. Then \mathcal{PT}_n *partially* acts on the left of L

$$h * f = \begin{cases} hf & hf \in L \\ \text{undefined} & hf \notin L \end{cases} \quad (f \in L, \quad h \in \mathcal{PT}_n)$$



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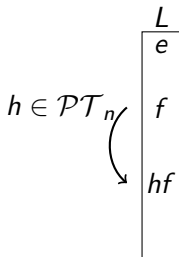
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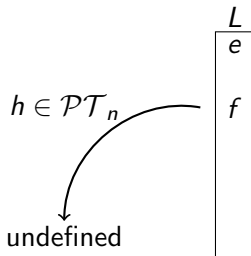
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Induced left Schützenberger module

- This induces a structure of $\mathbb{C}\mathcal{PT}_n - \mathbb{C}S_k$ bi-module on $\mathbb{C}L = \{\sum \alpha_i f_i \mid \alpha_i \in \mathbb{C}, f_i \in L\}$ defined on basis elements by

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- Let S^λ be a simple module of $\mathbb{C}S_k$. So $\mathbb{C}L \otimes S^\lambda$ is a $\mathbb{C}\mathcal{PT}_n$ -module.

Theorem (Munn, Ponizovskii (1950~), Margolis & Steinberg (2011)-modern proof)

$\mathbb{C}L \otimes S^\lambda$ is a simple $\mathbb{C}\mathcal{PT}_n$ -module and every simple module is of this form (for some choice of \mathcal{L} -class L and $\lambda \vdash k$).

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Remark

$$\mathbb{C}L_1 \otimes S^{\lambda_1} \simeq \mathbb{C}L_2 \otimes S^{\lambda_2} \iff \text{rank}(L_1) = \text{rank}(L_2), \quad \lambda_1 = \lambda_2$$

Construction of ind. projective modules of \mathcal{PT}_n

- Partial functions $f, g \in \mathcal{PT}_n$ are $\tilde{\mathcal{L}}_E$ -related $\iff \text{dom}(f) = \text{dom}(g)$.
Clearly $\mathcal{L} \subseteq \tilde{\mathcal{L}}_E$.
- For every $A \subseteq \{1, \dots, n\}$ denote by id_A the partial identity of A .

$$\text{id}_A(i) = \begin{cases} i & i \in A \\ \text{undefined} & i \notin A \end{cases}$$

- Note that every $\tilde{\mathcal{L}}_E$ class contains one partial identity.

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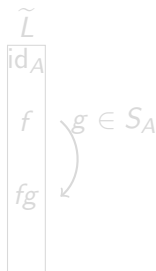
$$\text{id}_A(i) = \begin{cases} i & i \in A \\ \text{undefined} & i \notin A \end{cases}$$

- Note that every $\tilde{\mathcal{L}}_E$ class contains one partial identity.

Natural actions on an $\tilde{\mathcal{L}}_E$ -class

Let \tilde{L} be an $\tilde{\mathcal{L}}_E$ -class and let $\text{id}_A \in \tilde{L}$. Then S_A acts on the right of \tilde{L} by right multiplication

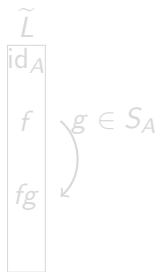
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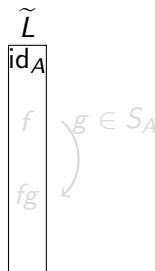
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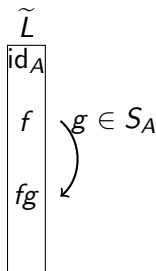
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Let \tilde{L} be an $\tilde{\mathcal{L}}_E$ -class. Then \mathcal{PT}_n *partially* acts on the left of \tilde{L}

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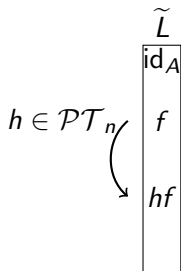
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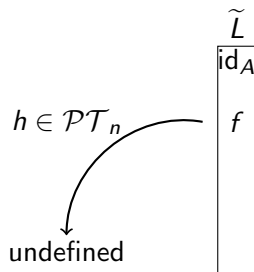
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Ind. projective modules

- This induces a structure of $\mathbb{C}\mathcal{PT}_n - \mathbb{C}S_A$ bi-module on $\mathbb{C}\tilde{L} = \{\sum \alpha_i f_i \mid \alpha_i \in \mathbb{C}, f_i \in \tilde{L}\}$ as before.
- Let S^λ be a simple module of $\mathbb{C}S_A$. So $\mathbb{C}\tilde{L} \otimes S^\lambda$ is a $\mathbb{C}\mathcal{PT}_n$ -module.

Theorem (Margolis & IS)

$\mathbb{C}\tilde{L} \otimes S^\lambda$ is an ind. projective $\mathbb{C}\mathcal{PT}_n$ -module and every ind. projective module is of this form (for some choice of \tilde{L}_E -class \tilde{L} and $\lambda \vdash |A|$).

Remark

$$\mathbb{C}\tilde{L}_1 \otimes S^{\lambda_1} \simeq \mathbb{C}\tilde{L}_2 \otimes S^{\lambda_2} \iff |\text{dom}(\tilde{L}_1)| = |\text{dom}(\tilde{L}_2)|, \quad \lambda_1 = \lambda_2$$

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Dimension of Ind. Projective modules

If $|\text{dom}(\tilde{L})| = k$ and $\lambda \vdash k$ then

$$\dim(\mathbb{C}\tilde{L} \otimes S^\lambda) = \sum_{l=1}^k \sum_{\substack{\mu \vdash k \\ |\mu|=l}} \binom{n}{l} K_{\lambda\mu}$$

(where $K_{\lambda\mu}$ is the Kostka number of λ and μ - the number of semistandard Young tableaux with shape λ and content μ).

New? - category approach

- Define a category \mathcal{E}_n as follows. Objects: subsets of $\{1, \dots, n\}$. Morphisms: $\mathcal{E}_n(A, B) = \{f : A \rightarrow B \mid f \text{ onto}\}$

Theorem (IS 2016)

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- Similar to other results in the literature (bands, \mathcal{R} -trivial monoids etc.).
- Transparent correspondence between simple and its projective cover. Let L and \tilde{L} be the \mathcal{L} and $\tilde{\mathcal{L}}$ -classes of some partial identity id_A . Define

$$\psi : \mathbb{C}\tilde{L} \rightarrow \mathbb{C}L$$

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General case

Everything said can be generalized to a class of **finite** semigroups satisfying three properties:

- 1 E -Ehresmann
- 2 right restriction
- 3 the corresponding Ehresmann category is EI (every endomorphism is an isomorphism)

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Ehresmann semigroups

Definition (Lawson, 1991)

Let S be a semigroup and $E \subseteq S$ be a subsemilattice. S is called E -Ehresmann if the following holds:

- For every $a \in S$ there is a minimal $e \in E$ such that $ea = a$ (denoted a^+).
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- The following identities hold: $(ab)^+ = (ab^+)^+$ and $(ab)^* = (a^*b)^*$.

Example

$S = \mathcal{PT}_n$, $E = \{\text{id}_A \mid A \subseteq \{1, \dots, n\}\}$, $f^+ = \text{id}_{\text{im}(f)}$ and $f^* = \text{id}_{\text{dom}(f)}$.

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Ehresmann categories - categories Isomorphism

To every E -Ehresmann S semigroup we can associate an Ehresmann category \mathcal{C} . The set of objects is E and the morphisms are in one-to-one correspondence with elements of S .

$$a \in S \implies C(a) : a^* \rightarrow a^+$$

$$C(b) \cdot C(a) = \begin{cases} C(ba) & b^* = a^+ \\ \text{undefined} & \text{otherwise} \end{cases}$$

We require this category to be an EI - category.

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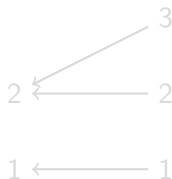
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Right restriction

An E -Ehresmann semigroup S is *right restriction* if the following identity holds:

$$ea = a(ea)^*, \quad (e \in E, \quad a \in S)$$

- For $S = \mathcal{PT}_n$, $f \text{id}_A$ is restriction of the domain to A and $\text{id}_A f$ is restriction of the image to A . In \mathcal{PT}_n , an image restriction can be achieved by domain restriction but not vice versa!



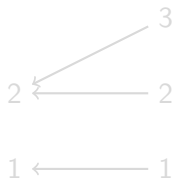
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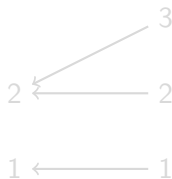
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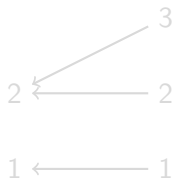
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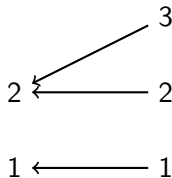
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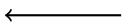
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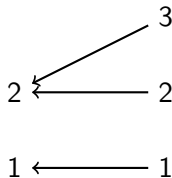
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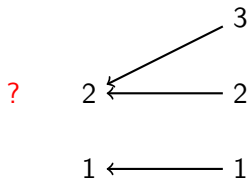
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Let S be a finite, E -Ehresmann and right restriction semigroup with corresponding EI-Ehresmann category. Let $a, b \in S$ then $a \mathcal{L} b$ if they generate the same left principal ideal: $S^1 a = S^1 b$. Let L be a regular \mathcal{L} -class (contains an idempotent $e \in L$). Then we can construct as above an $\mathbb{C}S$ -module $\mathbb{C}L \otimes V$ (where V is a simple $\mathbb{C}G_e$ -module).

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$\mathbb{C}L \otimes V$ is a simple $\mathbb{C}S$ -module and every simple module is of this form.

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Ind. Projective modules

Let S be a finite, E -Ehresmann and right restriction semigroup with corresponding EI-Ehresmann category. Let $a, b \in S$ then $a \tilde{\mathcal{L}}_E b$ if $a^* = b^*$. Let \tilde{L} be an $\tilde{\mathcal{L}}_E$ -class (which contains precisely one idempotent $e \in E$). Then we can construct as above an $\mathbb{C}S$ -module $\mathbb{C}\tilde{L} \otimes V$ (where V is a simple $\mathbb{C}G_e$ -module).

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Final remarks

- The class of all finite E -Ehresmann right restriction semigroup with corresponding EI-Ehresmann category is a pseudovariety \mathcal{V} of bi-unary semigroups.
(pseudovariety - closed under subalgebra, homomorphic image and *finite* direct products).
- Clearly, every $(2,1,1)$ -subalgebra of \mathcal{PT}_n is in \mathcal{V} .

Question

Is every element of \mathcal{V} a $(2,1,1)$ -subalgebra of \mathcal{PT}_n ?

Proposition (Margolis & IS)

If $S \in \mathcal{V}$ and *regular* then it is embeddable in \mathcal{PT}_n as a bi-unary semigroup.

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Thank you!